

Property-Based Test-Case Generators for Free

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Property-Based Testing (PBT)

Idea behind PBT

instead of designing specific test-cases,
a software test engineer **specifies properties**
of the **program inputs and outputs**

Then, a test execution engine

1. randomly generates an **input** that **satisfies** the specification
2. runs the program with that input
3. checks whether or not the **output** **satisfies** the specification

QuickCheck (Haskell, [Classen & Hughes, ICFP '00])

PropEr (Erlang, [Papadakis et al., ACM SIGPLAN WKSH Erlang '11])

Constraint-based PBT

FocalTest [Carlier et al., ICSOFT '10, TAP '12]

PBT for Focalize based on Constraint Logic Programming

other **constraint-based** testing methods

ArbitCheck (Java), JML-TT (Java/JML), PathCrawler (C), DART (C),
CUTE (C), Euclide (C), PrologCheck (Prolog), GATeL (Lustre),
CutEr (Erlang), AUTOFOCUS (MBT XP dev. Proc.), ...

This work

ProSyT: Property-based Symbolic Testing

PBT framework for testing **Erlang** programs
based on Constraint Logic Programming

smart symbolic execution of an **interpreter** for Erlang
to generate input test values
directly from program inputs specifications

Property-Based Testing for Erlang

Erlang: concurrent, higher-order, functional programming language with dynamic, strong typing

Erlang program: sequence of function definitions

$$f(X_1, \dots, X_n) \rightarrow e$$

f is a function name, X_1, \dots, X_n are variables, and **e** is an expression

A **faulty**
insertion
program

```
insert(E,L) ->
  case L of
    [] -> [E];
    [X|Xs] when E=<X -> [X,E|Xs];
    [X|Xs] -> [X] ++ insert(E,Xs)
  end.
```

Erlang PBT: PropEr (Property-based testing tool for Erlang)

<https://github.com/proper-testing/proper>

Specifying PBT tasks with PropEr

Property to be satisfied by the program inputs and outputs

Given any integer I and any ordered list of integers L,
insert(I,L) produces an ordered list

prop_ordered_insert() ->

?**FORALL**({I,L},

{integer(),ordered_list()},

ordered(insert(I,L))).

Specification of the
GENERATOR
of program inputs

Specification of the
PROPERTY
of program outputs

Properties of program outputs

```
prop_ordered_insert() ->
  ?FORALL( {I,L},
           {integer(),ordered_list()},
           ordered(insert(I,L)) ).
```

An Erlang boolean function:

```
ordered(L) -> case L of
  [A,B|T] -> A =< B andalso ordered([B|T]);
  _ -> true
end.
```

Properties of program inputs

```
prop_ordered_insert() ->  
  ?FORALL( {I,L},  
           {integer(),ordered_list()} ,  
           ordered(insert(I,L)) ).
```

Generators are Erlang expressions built upon

- Predefined types

`integer()`, `float()`, `list(integer())`, `{ integer(), ... }`

- User-defined types

`-type tree() :: 'leaf' | {'node',tree(T),T,tree(T)}`

`ordered_list()` ->

`?SUCHTHAT(L, list(integer()), ordered(L))`.

Filter of valid lists

Specifying filters

```
prop_ordered_insert() ->  
  ?FORALL( {I,L},  
           {integer(),ordered_list()},  
           ordered(insert(I,L))      ).
```

- **without an** ad-hoc implementations of `ordered_list()`, the generation of **ordered lists** from

```
ordered_list() ->  
  ?SUCHTHAT(L, list(integer()), ordered(L)).
```

is performed in an **inefficient way** by randomly **generating** a list of integers & **testing** if ordered

- implementing ad-hoc generators **time-consuming** & **error-prone** activity

Our goal & contribution

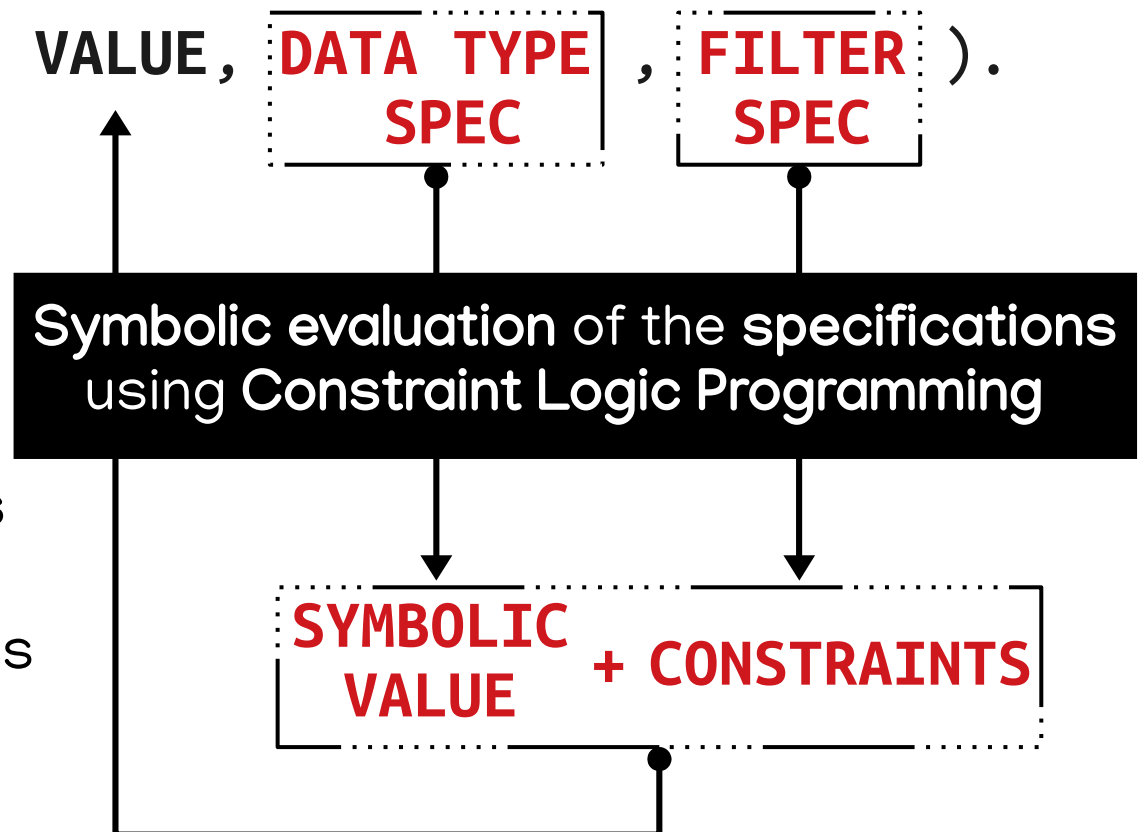
To relieve developers from implementing ad-hoc input generators by deriving generators directly from **PropEr specifications**

ordered_list() ->

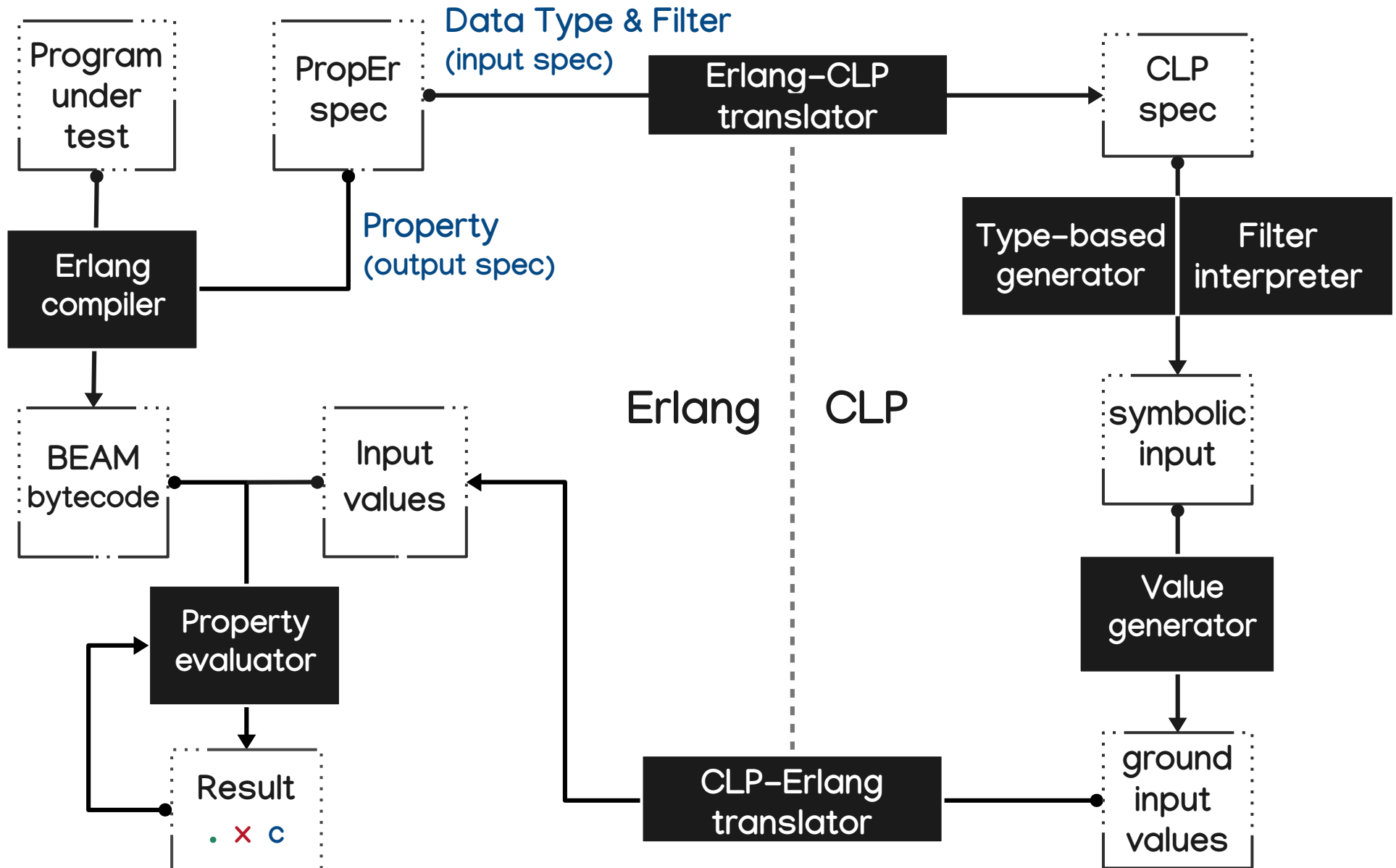
?SUCHTHAT(VALUE, DATA TYPE SPEC, FILTER SPEC).

Constrain & generate computation pattern:

1. generates **symbolic data structures** with **constraints**
2. generates random values satisfying those constraints
3. translates back concrete data into inputs for the program under test



Property-Based Symbolic Testing



Constraint Logic Programming

CLP(X) languages

CLP(FD) for integers

CLP(R) for floats

CLP(B) for booleans

Constraint of CLP(FD)

Quantifier free first-order formula
whose variables range over FD

Definition of the predicate `insert/3`:

clauses $\left\{ \begin{array}{l} \text{insert}(I, [], [I]). \\ \text{insert}(I, [X|Xs], [I, X|Xs]) \text{ :- } I \#=< X. \\ \text{insert}(I, [X|Xs], [X|Ys]) \text{ :- } I \#> X, \text{insert}(I, Xs, Ys). \end{array} \right.$

query \longrightarrow `?- I #>= 3, I #=< 7, L = [2,4,8], insert(I,L,0).`

`L = [2,4,8], 0 = [2,I,4,8], I in 3..4 ;`
`L = [2,4,8], 0 = [2,4,I,8], I in 5..7 ;` } answers
`false.`



SWI Prolog

Translator from PropEr to CLP

```
prop_ordered_insert() ->
```

```
  ?FORALL( {I,L},  
           {integer(),ordered_list()},  
           ordered(insert(I,L)) ).
```

```
ordered_list() ->
```

```
  ?SUCHTHAT( L, list(integer()), ordered(L) ).
```

Erlang

CLP

```
prop_ordered_insert_input(I,L) :-  
  typeof(I,integer), ordered_list(L).
```

```
ordered_list(L) :-
```

```
  typeof(L,list(integer)), eval(apply('ordered',[var('L')]),  
                                [('L',L)],  
                                lit(atom,true) ).
```

CLP translation of the

- data type
- filter function specifications

Type-based value generator

Given a data type specification T ,
the predicate `typeof(X,T)` holds iff
 X is a CLP term encoding an Erlang value of type T .

```
ordered_list() ->  
    ?SUCHTHAT(L, list(integer()), ordered(L)).
```

```
typeof(nil,list(T)).  
typeof(cons(Hd,Tl),list(T)) :-  
    typeof(Hd,T), typeof(Tl,list(T)).
```

```
?- typeof(L,list(integer)).  
L = nil ;  
L = cons(lit(int,X),nil), X in inf..sup ;  
...
```

Size (of terms)
is configurable:

- length of lists
- interval of integers
- ...

Interpreter of filter functions

The CLP interpreter provides the predicate

`eval(In, Env, Out)`

that computes

the **output** expression `Out` from

the **input** expression `In` in the **environment** `Env`

(maps variables to values)

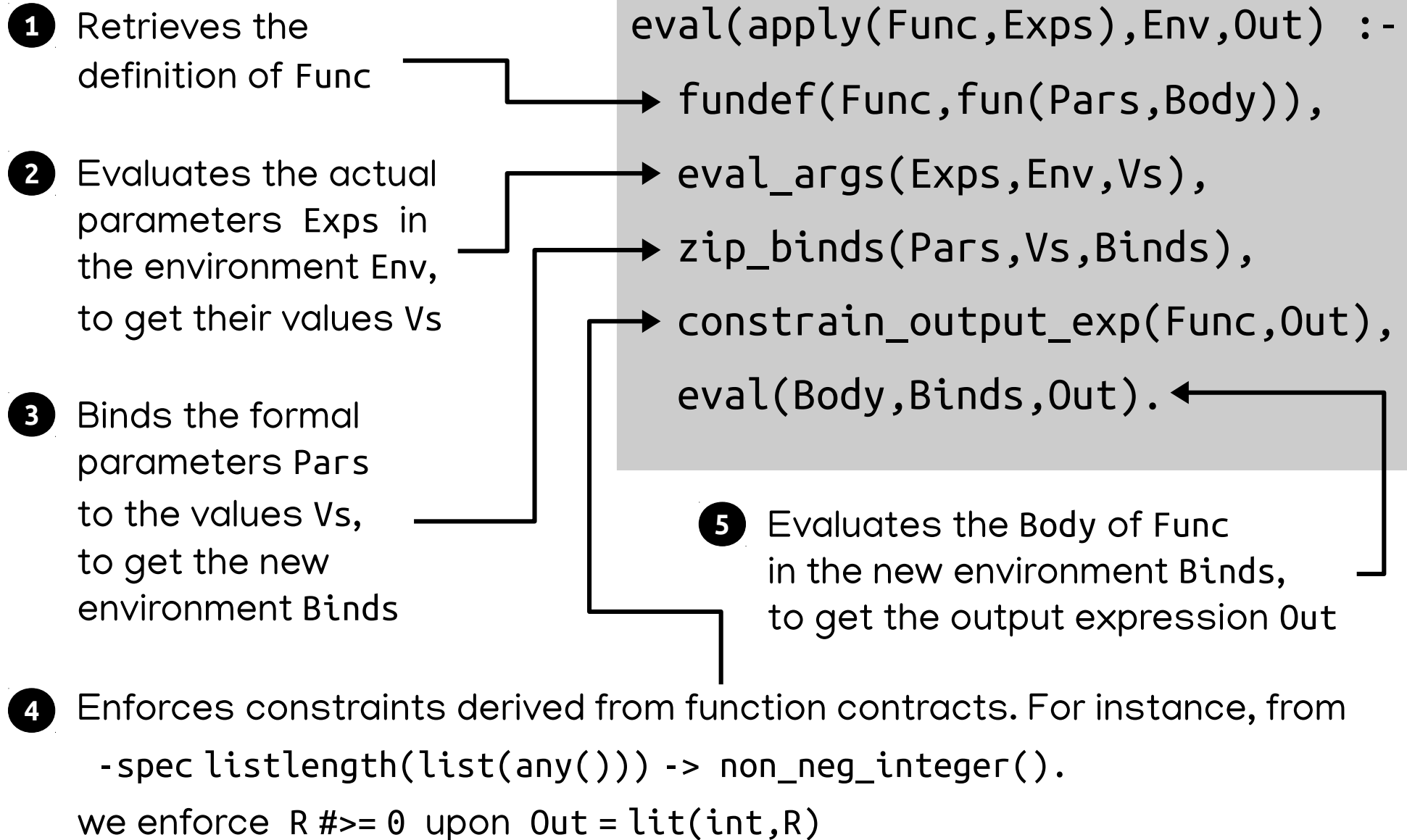
Using **symbolic** expressions (`In` and `Out` are CLP terms possibly with variables) in `eval` enables the

exploration of all program computations

without

explicitly enumerating all concrete inputs

Interpreter of filter functions



Generating symbolic ordered lists

```
ordered_list(L) :-                                     generetes a non-ground CLP
  typeof(L,list(integer)),                             term representing a list
  eval(apply('ordered',[var('L')]),[('L',L)],lit(atom,true)).
```

enforces constraints on the list
(ascending order of its elements)

```
?- ordered_list(L).
L = nil ;
L = cons(lit(int,X),nil), X in inf..sup ;
L = cons(lit(int,X),cons(lit(int,Y),nil)),
  Y #>= X ;
L = cons(lit(int,X),cons(lit(int,Y),cons(lit(int,Z),nil))),
  Y #>= X, Z #>= Y
```


Value generator

Random generation of **ground** terms

```
rand_elem(nil).
rand_elem(cons(X,L)) :- rand_elem(X), rand_elem(L).
rand_elem(lit(int,V)) :-
    fd_inf(V,Inf), fd_sup(V,Sup), random_between(Inf,Sup,V).
```

```
?- ordered_list(L), rand_elem(L), write_elem(L).
```

```
[]
```

```
L = nil ;
```

```
[10]
```

```
L = cons(lit(int,10),nil) ;
```

```
[4,8]
```

```
L = cons(lit(int,4),cons(lit(int,8),nil)) ;
```

```
[2,6,9]
```

```
L = cons(lit(int,2),cons(lit(int,6),cons(lit(int,9),nil)))
```

Translate
CLP terms to
Erlang values

Running ProSyT

<https://fmlab.unich.it/testing/>

```
$ ./prosynt.sh ord_insert_bug.erl prop_ordered_list \
```

```
--min-size 10 --max-size 100
```

```
--inf -1000 --sup 1000
```

```
--tests 250 --verbose
```

size of the data structure

interval where the integers are taken from

number of tests to run

Tests Results:

```
XXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXX  
XXXXXXXXXXXXXXXXXXXXX.XXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXX  
XXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXX  
XXXXXXXXXXXXXXXXXXXXX.XXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXX  
XX
```

Timings (s)

erl2clp		0.58
tests generation		0.12 (test cases generated: 250 out of 250)
testing		3.89

Generating complex data structures

Symbolic test-case generation performs well when the filter **does not specify constraints on the skeleton** of the data structure, but only on its elements

AVL tree

- binary search tree (constraints on the elements)
- height-balanced (constraints on the skeleton)

```
avl(T) -> case T of
  leaf -> true;
  {node,L,V,R} -> B = height(L) - height(R)      andalso
                  B >= -1      andalso      B =< 1      andalso
                  ltt(L,V)      andalso      gtt(R,V)      andalso
                  avl(L)      andalso      avl(R)      ;
  _ -> false
end.
```

Symbolic generation of AVL trees

```
?- typeof(X,tree(integer)),  
   eval(apply('avl',[var('T')]),[( 'T',X)],lit(atom,true)).
```

- ① generates a symbolic **binary tree** X
- ② applies the filter to X
 - makes X a **search tree**
by enforcing **constraints on the values** of the nodes
 - makes X **height-balanced**
how? Can't enforce **constraints on the skeleton** of X
which is determined by ①

Among the answers of ① just a few are height-balanced trees

For trees of size 10 (number of nodes) ② finds 10 AVL trees
out of 9000 binary trees generated by ①

Data-driven generation: coroutining

Interleaving the execution of

- the type-based generator `typeof`, and
- the interpreter of filter functions `eval`

enables **enforcing constraints while generating** data

```
eval(apply('avl', [var('T')]),  
      [('T', X)], lit(atom, true)), typeof(X, tree(integer))
```



`typeof` and `eval` cooperate through `X`
during the generation of the AVL tree

Implemented using the **coroutining** mechanism provided by SWI-Prolog through the primitive

```
when(Cond, Conj)
```

that **suspends** the execution of `Conj` until `Cond` becomes true

Coroutining typeof and eval

```
?- typeof(...), eval(...).
```

```
eval(case(CExps,Cls),Env,Exp) :-  
  eval(CExps,Env,EExps),  
  suspend_on(Env,EExps,Cls,Cond),  
  when(Cond, (  
    match(Env,Eexps,Cls,MEnv,C1),  
    eval(C1,MEnv,Exp)  
  )).  
}
```

selects the variables that would get bound to **lists** or **tuples** while matching EExpr against the clauses Cls of the case-of expression

suspends the evaluation of the match until all the variables get bound non-variable terms

```
?- eval(...) corouting, typeof(...).
```

Generating AVL trees

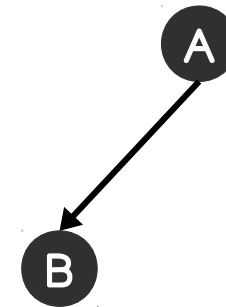
```
?- eval(apply('avl',[var('T')]),  
      [ ('T',X)],lit(atom,true), coroutingtypeof(X,tree(integer))
```

height-balanced (BST) ?

The evaluation of the filter adds constraints on X

$$\begin{array}{l} \underbrace{\hspace{10em}}_1 \\ -1 \#=< \text{height}(A_{\text{left}}) - \text{height}(A_{\text{right}}) \#=< 1, \\ -1 \#=< \text{height}(B_{\text{left}}) - \text{height}(B_{\text{right}}) \#=< 1, \\ A \#< B \\ \underbrace{\hspace{10em}}_0 \end{array}$$

the constraints on X restrict the possible ways in which its left and right subtrees can be further expanded



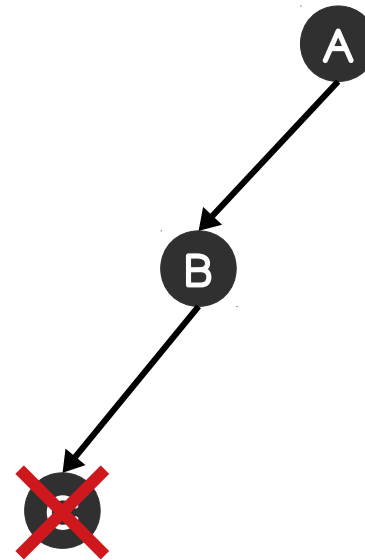
As soon as the typeof (partially) instantiates X to a binary tree

Generating AVL trees

```
?- eval(apply('avl',[var('T')]),  
        [ ('T',X)],lit(atom,true), coroutingtypeof(X,tree(integer)))
```

height-balanced (BST) ?

$-1 \#=< \overbrace{\text{height}(A_{\text{left}}) - \text{height}(A_{\text{right}})}^2 \text{ unsat} \#=< 1,$
 $-1 \#=< \underbrace{\text{height}(B_{\text{left}}) - \text{height}(B_{\text{right}})}_1 \#=< 1,$
 $A \#< B$

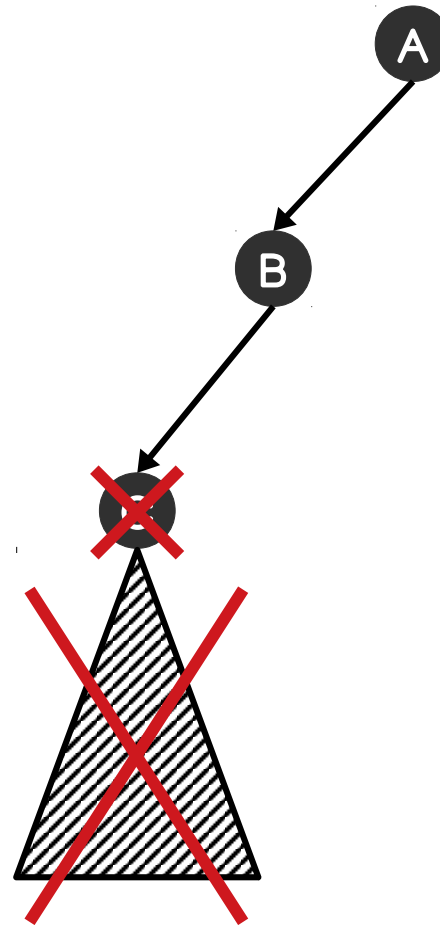


Generating AVL trees

```
?- eval(apply('avl',[var('T')]),  
        [ ('T',X)],lit(atom,true), coroutingtypeof(X,tree(integer)))
```

height-balanced (BST) ?

$-1 \#=< \overbrace{\text{height}(A_{\text{left}}) - \text{height}(A_{\text{right}})}^2 \text{ unsat} \#=< 1,$
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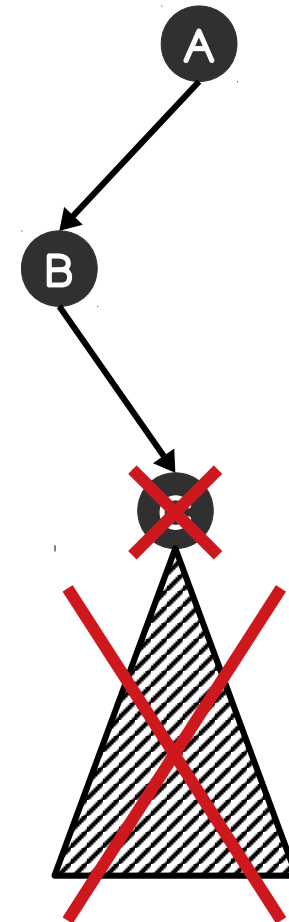


Generating AVL trees

```
?- eval(apply('avl',[var('T')]),  
        [ ('T',X)],lit(atom,true), coroutingtypeof(X,tree(integer))
```

height-balanced (BST) ?

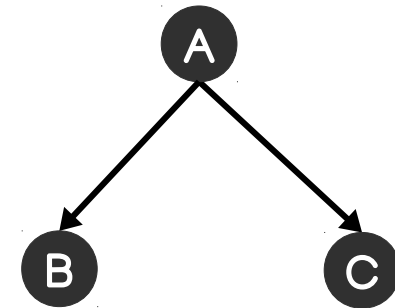
$-1 \#=< \overbrace{\text{height}(A_{\text{left}}) - \text{height}(A_{\text{right}})}^2 \text{ unsat} \#=< 1,$
 $-1 \#=< \underbrace{\text{height}(B_{\text{left}}) - \text{height}(B_{\text{right}})}_{-1} \#=< 1,$
 $A \#< B$



Generating AVL trees

```
?- eval(apply('avl',[var('T')]),  
        [ ('T',X)],lit(atom,true), coroutingtypeof(X,tree(integer)))
```

height-balanced (BST) ?

$$\begin{array}{l} \phantom{-1 \#=<} \phantom{\text{height}(A_{\text{left}})-\text{height}(A_{\text{right}})} \phantom{\#=<} \\ \phantom{-1 \#=<} \phantom{\text{height}(A_{\text{left}})-\text{height}(A_{\text{right}})} \phantom{\#=<} \\ -1 \#=< \underbrace{\text{height}(A_{\text{left}})-\text{height}(A_{\text{right}})}_{0} \#=< 1, \\ -1 \#=< \underbrace{\text{height}(B_{\text{left}})-\text{height}(B_{\text{right}})}_{0} \#=< 1, \\ A \#< B \end{array}$$


Experimental evaluation

Program	PropEr		ProSyT	
	Time	N	Time	N
ord_insert	300.00	0	300.00	67,083
up_down_seq	300.00	0	300.00	22,500
n_up_seqs	300.00	0		24,000
delete	300.00	0	9.21	100,000
stack	143.71	100,000	19.57	100,000
matrix_mult	300.00	0	300.00	76,810
det_tri_matrix	300.00	304	32.28	13,500
balanced_tree	300.00	121	21.54	100,000
binomial_tree_heap	300.00	0	43.45	4,500
avl_insert	300.00	0	300.00	23,034

Time reports the seconds needed to generate N ($\leq 100,000$) test cases of size in the interval $[10, 100]$ within the time limit of 300s. (Intel® Core™ i7-8550U with 16GB of memory running Ubuntu 18.04.2 LTS)

Conclusions

ProSyT

a PBT framework that relieves developers from writing generators of input values for testing Erlang programs

- ✓ based on a **constrain & generate** computation pattern implemented using a **CLP interpreter** that makes the generation process efficient in many cases
- ✓ CLP is fully **transparent** to users
 - translator from PropEr/Erlang specifications to CLP
 - translator from CLP test-cases to Erlang

Future work

- provide developers with explicit **shrinking** mechanism
- apply the approach to other programming languages, the **interpreter** makes it independent of the prog. lang.